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For: METHOD TO AUTHENTICATE A MOBILE STATION, A COMMUNICATIONS SYSTEM  
AND A MOBILE STATION

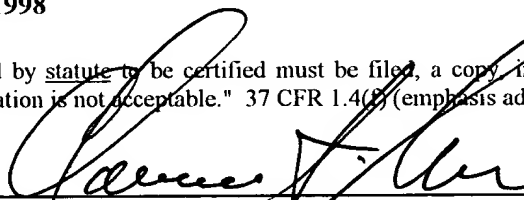
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"Method to authenticate a mobile station, a communications system and a mobile station"

(Menetelmä matkaviestimen todentamiseksi, tietoliikennejärjestelmä ja matkaviestin)

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## Method to authenticate a mobile station, a communications system and a mobile station

5 The target of the invention is the method to authenticate a mobile station specified in the preamble of the claim 1, the communications system specified in the preamble of the claim 9 and the mobile station specified in the preamble of the claim 11.

10 Known authentication and key agreement protocols are based either on symmetric or public key cryptography and a trusted third party. In GSM the authentication and encryption key agreement is based on symmetric key and a trusted third party. The method using symmetric key require the existence of an agreed secret between communicating parties or with a server as the third party. In GSM the mobile station of the subscriber shares a secret subscriber authentication key  $K_i$  with a trusted authentication centre AC. The authentication of the mobile station is based on the use of a one-way function A3 and a ciphering key  $K_c$  is derived from the shared  $K_i$  in 15 the mobile station and the authentication centre.

Figure 1 shows a prior art authentication arrangement of GSM mobile networks, where there is an Authentication Centre AC 1, a Home Location Register HLR 2, Visitor Location Register VLR 3, Base Transmitter Station BTS 4 and Mobile 20 Equipment ME 5, where number 6 is a Subscriber Identity Module SIM.

The method by the arrangement operates as follows: Authentication Centre 1 forms a Random Number RAND, that is used with subscriber authentication key  $K_i$  to form an authentication triplet 7. The authentication triplet 7 comprises random number RAND directly from the RAND above, Signed Response SRES formed with a one-way function A3 1a from the subscriber authentication key  $K_i$  and 25 ciphering key  $K_c$  formed with one-way function A8 1b from the RAND above. The authentication triplet 7 is sent to Home Location Register HLR 2 and then to Visitor Location Register VLR 3. The RAND of the authentication triplet 7 is sent from the VLR 3 to the Subscriber Identity Module SIM 6 in the Mobile Equipment ME 5 to form a key corresponding to the ciphering key  $K_c$  in the same authentication triplet 7. The above key is formed by one-way function A8 6b in SIM 6 and processed more by one-way function A5 8 in ME 5 to exchange with the  $K_c$  of the authentication triplet 7 processed by one-way function A5 8 in the Base Transmitter 30

Station 4. Also subscriber authentication key  $K_i$  of the SIM 6 is used to form a signed response corresponding to the SRES in the above authentication triplet 7 in the VLR 3. This signed response is directly sent to the VLR 3 to compare it with the SRES to complete the authentication.

- 5 Formerly is also known User-to-User Signalling (UUS) that is defined for Integrated Services Digital Network (ISDN) and is being defined for GSM network. The UUS is defined for GSM in ETSI (European Telecommunications Standards Institute) specification Digital cellular telecommunications system (Phase 2+); User-to-User Signalling (UUS); Service description, Stage 1 (GSM 02.87).
- 10 The UUS supplementary service allows the served subscriber to send to or receive from another user a limited amount of information. This information is generated by the subscriber and shall be passed transparently through the network. With the word transparently is meant that no modification to the contents is made. Normally the network does not interpret this information.
- 15 The served subscriber is able to send and receive User-to-User Information (UUI) in different phases of the call depending on what service subscriber uses. Possible services are:
  - 20 Service 1: UUI can be sent and received during the origination and termination of a call, with UUI embedded within call control messages. The service 1 can be activated implicit by inserting UUI when set-up a call or explicit with an appropriate procedure.
  - 25 Service 2: UUI can be sent and received after the served subscriber has received an indication that the remote party is being informed of the call and prior to the establishment of the connection. UUI sent by the served subscriber prior to receiving the acceptance of the call by the remote party, may as a network option be delivered to the remote party after the call has been established. The service 2 shall be activated explicitly.
  - Service 3: UUI can be sent and received only while the connection is established. The service 3 shall be activated explicitly.
  - 30 Services 1 to 3 shall allow the transmission of UUI with the maximum length of 128 octets per message. In some networks as ISDN the maximum length is only 32

octets. The USER INFORMATION message between GSM mobile station and Mobile Switching Centre (MSC) can have 128 octets of user data while the messages for call setup and release can have 32 octets of user data. Messages for call setup and release include f.ex. SETUP, PROGRESS, ALERT, CONNECT, DISCONNECT, RELEASE, RELEASE COMPLETE.

- A problem in the known arrangements in mobile networks is that only security parameters required in the establishment of confidentiality on the air interface are exchanged. Parameters for other information security features as integrity on air interface are not agreed.
- 10 A problem in the shared-key authentication and key agreement procedures is the agreed secret and in some circumstances needed connection to the trusted third party during the execution of the protocol.

The objective of the invention is to avoid disadvantages of the prior art solutions by bringing out a new authentication and key agreement for mobile communications systems to implement an end-to-end secure transmission.

- What is characterizing to the method according to the invention is presented in claim 1. What is characterizing to the communications system of the invention is presented in claim 9. What is characterizing to the mobile station of the invention is presented in claim 11. Preferable embodiments of the invention are described in the dependent claims.

The invention concerns a method to authenticate a mobile station in a mobile network. According to the invention the mobile station is authenticated using user-to-user data exchange. This can be done during call setup or call.

- In one embodiment of the method also an encryption key is agreed between mobile stations.

The invention concerns also a cellular communications system, where the first and second mobile stations (A, B) are connected wireless with via base stations. According to the invention the cellular communications system comprises a first mobile station (A), that constructs and sends a first message ( $M_1$ ), receives and verifies the validity of a second message ( $M_2$ ) and when the information is verified valid accepts to share a shared encryption key  $K$ , constructs and sends a third message

(M<sub>3</sub>), a second mobile station (B), that receives the first message (M<sub>1</sub>) and constructs and sends the second message (M<sub>2</sub>), receives and verifies the validity of the third message (M<sub>3</sub>) and when the information is valid accepts to share the shared encryption key K with the first mobile station (A), and at least one mobile switching centre.

The invention concerns also a mobile station. According to the invention the mobile station comprises a processor to perform operations needed to form and verify messages (M<sub>1</sub>, M<sub>2</sub>, M<sub>3</sub>), to implement authentication and key agreement procedures, a memory, where procedures and messages are stored with necessary parameters and variables, output means, on which commencement of extra secure communication is presented to a user of the mobile station, input means to enable validation of the extra secure communication, a transmitter/receiver and an antenna to transform information to radio waves from digital signals and vice versa.

An advantage of the invention is that the traffic between the communicating mobile stations is protected autonomously with public-key based authentication and key agreement mechanisms.

The invention is described in detail in the following by referring to the appending drawing, where

figure 1 presents a prior art arrangement in a flow chart,

figure 2 presents a method of the invention in a signalling diagram,

figure 3 presents a signalling diagram of an authentication and key agreement protocol,

figure 4 presents communication system of the invention, and

figure 5 presents essential parts of mobile station according to the invention in a block diagram.

Figure 1 is described in the prior art portion of the text.

Figure 2 shows an authentication and key agreement protocol. The protocol is started when an input is given to trigger extra secure transmission. The parenthesis

after the name of the message contain the name of the part of the message where the carried information is included and the name of the information. First the calling mobile station (MS) 9 sends a SETUP(UUS(service code)) message to the first Mobile Switching Centre (MSC) 10. The SETUP message contains in a User-to-  
 5 User Information (UUI) element a service code indicating the encryption key management service encoded by the calling mobile station 9. The UUI element is transferred with the User-to-User Signalling (UUS). The first MSC 10 sends the user-to-user information via the Integrated Service Digital Network (ISDN) User Part (ISUP) signalling in an Initial Address Message (IAM) (UUS(service code)) to  
 10 the second MSC 11 to which the called mobile station 12 is connected. This signalling between two MSCs 10, 11 is only needed when the two mobile stations 9, 12 are connected to different MSCs 10,11. The first MSC 10 responds to the calling mobile station 9 with a CALL PROceeding message and the second MSC 11 sends a SETUP(UUS(service code)) message formed by the data from the first SETUP  
 15 message to the called mobile station 12. Now the mobile station informs the user about an extra secure call. The called mobile station 12 responds to the second MSC 11 with a CALL CONFirmed and an ALERT(UUS(service acc.)) message meaning that the terminal equipment is alerting the subscribed user. Information whether the called mobile station 12 accepts the extra secure communication is delivered in the  
 20 ALERT message. The ALERT message is led to the calling mobile station 9 to inform the alerting and the possible acceptance. The information is transferred if needed between two MSCs 11, 10 in an ISUP Answer Message (ANM). In case the extra secure communication is not applied preferably a normal call setup is continued or the call setup is aborted. This can be commenced by a decision of the user  
 25 or of the logic of the mobile station 9 or of the logic of the MSC 10.

If the extra secure communication is accepted the service, the authentication and the key agreement protocol related information is exchanged between the two mobile stations 9, 12 in the USER INFOrmation message of GSM and ISUP. First the USER INFO(UUS(M<sub>1</sub>)) message is transferred from the calling mobile station 9  
 30 through MSCs 10, 11 to the called mobile station 12. Then the USER INFO(UUS(M<sub>2</sub>)) message is transferred from the called to the caller and the USER INFO(UUS(M<sub>3</sub>)) message is transferred from the caller to the called. If one or more of the messages M<sub>1</sub>, M<sub>2</sub>, M<sub>3</sub> is longer than the space in one USER INFO message carrier several USER INFO messages are used for transportation.

35 At last, during the call setup the called mobile station 12 sends a CONNECT(UUS(data)) message to the MSC 11 it is connected with. And the MSC

11 responds with a CONNECT ACKnowledgement message to the called mobile station 12. Then if needed the MSC 11 sends an AMN(UUS(data)) message to another MSC 10. The MSC 10 connected with the calling mobile station 9 sends a CONNECT(UUS(data)) message to the calling mobile station 9 and receives a  
 5 CONNECT ACKnowledgement message. If User-to-User Signalling data UUS(data) is not needed at this stage plain CONNECT and AMN messages or messages with empty UUS(data) fields are used. It is possible to transfer more User-to-User Information now after the call is connected. The encryption algorithms can be applied to this information.

10 Alternatively, the security parameters can be exchanged after call setup during the call. In this option the call can be setup normally. When either subscriber wants to start extra secure communication during the conversation or data exchange, the subscriber initiates the secure communications e.g. by pressing the keys of the keyboard and the security parameters are exchanged using User-to-User Signalling.

15 Figure 3 shows a signalling diagram of one authentication and key agreement protocol that can be used in the inventive method. The messages  $M_1$ ,  $M_2$ ,  $M_3$  are shown without reference to the User-to-User Signalling (UUS) that the transmission is based on. The references A and B cite to the mobile stations in this station-to-station protocol. The references 1. to 6. cite to the steps performed while progressing.  
 20 Subscribers A and B need an agreement on a key for extra secure connection. The protocol works followingly:

In step 1 the subscriber A initiates the protocol and selects a prime number  $p$ , a generator  $a$  of the multiplicative group of integers modulo  $p$  when  $p \geq a \geq 2$  and a random secret  $x$  when  $p-2 \geq x \geq 1$ . Then A constructs and sends to B the message  
 25  $M_1$  containing

$$a, p, a^x \bmod p.$$

In step 2 the subscriber B receives the message  $M_1$  and afterwards generates a secret  $y$  when  $p-2 \geq y \geq 1$  and computes a shared key  $K = (a^x)^y \bmod p$ . Then B signs the concatenation of exponentials  $\{a^y, a^x\}$  and encrypts the result  $S_B\{a^y, a^x\}$  with the  
 30 shared key leading to  $E_K(S_B\{a^y, a^x\})$ . B constructs and sends the message  $M_2$  to A in step 3 containing

$$a^y \bmod p, cert_B, E_K(S_B\{a^y, a^x\}).$$



Certificate  $cert_B$  in the message  $M_2$  contains the signature verification key of the subscriber B. The exact contents of the certificate may differ from the following minimum

$$cert_B = (B, p_B, a, p, S_T\{B, p_B, a, p\}),$$

- 5 where  $p_B$  is the public signature verification key of the subscriber B and  $S_T$  is the signature transformation of a trusted authority T whose public signature verification key is known by A and B.

- In step 4 the subscriber A receives the message  $M_2$  and afterwards computes the shared encryption key  $(a^y)^x \bmod p = (a^x)^y \bmod p = K$ . The validity of the certificate  
 10  $cert_B$  is checked by the subscriber A. When the certificate  $cert_B$  is valid the encrypted part  $E_K(S_B\{a^y, a^x\})$  of the message  $M_2$  is decrypted to receive  $S_B\{a^y, a^x\}$  and the signature  $S_B\{a^y, a^x\}$  is verified with the public signature verification key  $p_B$  of the subscriber B. If the signature is verified valid A accepts to share the shared encryption key K with B. If the signature is invalid the execution of the protocol is  
 15 cancelled by A.

In step 5 the subscriber A signs the concatenation of exponentials  $\{a^x, a^y\}$  and encrypts the result  $S_A\{a^x, a^y\}$  with the shared key leading to  $E_K(S_A\{a^x, a^y\})$ . A constructs and sends the message  $M_3$  to B in step 5 containing

$$cert_A, E_K(S_A\{a^x, a^y\}),$$

- 20 where  $cert_A$  includes corresponding information with  $cert_B$  of the subscriber A. The exact contents of the certificate  $cert_A$  may differ from the following minimum

$$cert_A = (B, p_A, a, p, S_T\{B, p_A, a, p\}),$$

- where  $p_A$  is the public signature verification key of the subscriber A and  $S_T$  is the signature transformation of a trusted authority T whose public signature verification  
 25 key is known by A and B.

In step 6 the subscriber B receives the message  $M_3$  and verifies the validity of the  $cert_A$ , decrypts  $E_A(S_A\{a^x, a^y\})$  and verifies the validity of the signature of  $S_A\{a^x, a^y\}$ . If all the signatures are valid B accepts sharing of K with A. If any of the signatures is invalid B cancels the execution of the protocol.

Also other public key based authentication and key agreement protocols than the above presented station-to-station protocol can be used.

In figure 4 a communications system according to the invention is shown. The mobile station 9 of the subscriber A is connected wireless to a base transmitter station BTS that is connected wired to a base station controller BSC and to a mobile switching centre MSC 10. The MSCs 10, 11 are connected with an ISDN network together. The MSC 11 and BTS are connected wired and the BTS is connected wireless to the mobile station 12 of the subscriber B. Here only the MSCs 10, 11 are shown to present the logic of the invention. In reality the BTSs and the BSCs are also present.

In figure 5 a block diagram of the essential hardware needed to implement a mobile station according to the invention is described. The processor 13 perform the operations needed to implement the authentication and key agreement procedures from the memory 14 where they are stored with necessary parameters and variables. The commencement of extra secure communication is presented on the display 15 to the user of the mobile station. The validation of the service is done by pushing keys on the keyboard 16 or by processor 13. The transmitter/receiver 17 and antenna 18 is used to transform the information transmitted on radio waves from digital signals and vice versa.

The following example is presented to explain details of the invention when there are two different mobile switching centres connected with ISDN network together. User-to-User Signalling is used to transfer messages for station-to-station authentication and key agreement protocol described above. First the calling mobile station 9 encodes a service code indicating the encryption key management service to the user-to-user information element of the SETUP message and the mobile station 9 sends the message to the Mobile Switching Centre (MSC) 10. Then the User-to-User Information (UII) is transferred using the ISDN User Part (ISUP) signalling to the MSC 11 where the mobile station of the called subscriber 12 is connected if the subscribers are connected to different MSCs 11. The UII is transferred to the mobile station 12 of the called subscriber in the SETUP message. If the extra secure communication service defined in the UII is recognized in the mobile station 12, the called subscriber is alerted preferably with a sound and textual or symbolical way of the service. The user have to allow or refuse the service in concern. Information of allowance or refusal of the service is transferred in an ALERT message from the mobile station 12 to the mobile switching centre 11.

The UUI is transferred to the mobile station 9 of the calling subscriber preferably in an ISUP Answer Message (ANM) between the MSCs and in the ALERT message on the GSM connection. The calling subscriber is informed of the allowance or refusal of the extra secure communication service. If the called subscriber allowed the use of the service the authentication and key agreement protocol related information is exchanged between the two mobile stations using the USER INFO messages of GSM and ISDN. If the service is refused the call will be setup normally without the extra secure feature or the call will be terminated by user input without further setting up. The messages  $M_1$ ,  $M_2$ ,  $M_3$  are then exchanged as described in detailed descriptions of figures 2 and 3 above. Additional UUI information is transferrable between the mobile stations when the call is set up.

The extra secure communication can be initiated at least in three different ways followingly: 1. The user press a key or gives a voice command or gives an activation code before dialling a call, 2. A call to a subscriber on a list is made, and 3. The user chooses the feature from a menu to be on or off for a longer time. When the call is made the ability of the called mobile station to execute the required procedures is checked.

The examples described above are based on the use of the station-to-station protocol. The UUS signalling mechanism can be used to transport messages related to any other public key authentication and key agreement mechanism. It is also possible to use the UUS signalling mechanism to transport the messages of the shared-key technology based key agreement mechanisms.

The method of the invention can be used also in other networks that have a signalling mechanism between terminal equipment.

## Claims

1. A method to authenticate a mobile station in a mobile network, characterized in that the mobile station is authenticated with user-to-user data exchange.
2. A method according to claim 1, characterized in that the data is exchanged during call setup.
3. A method according to claim 1, characterized in that the data is exchanged during a call.
4. A method according to claim 1, 2 or 3, characterized in that also an encryption key is agreed between two mobile stations.
5. A method according to claim 4, characterized in that the mobile stations execute a mutual authentication and key agreement protocol based on public-key cryptography.
6. A method according to claim 5, characterized in that a second mobile station (B) is authenticated by
  - a first mobile station (A) constructing and sending to the second mobile station (B) a first message ( $M_1$ ), the second mobile station (B) receiving the first message ( $M_1$ ), constructing and sending a second message ( $M_2$ ) to the first mobile station (A), the first mobile station (A) receiving the second message ( $M_2$ ), checking the validity of the information in the second message ( $M_2$ ), if the information is verified valid the first mobile station (A) accepting to share a shared encryption key K with the second mobile station (B), the first mobile station (A) constructing and sending a third message ( $M_3$ ) to the second mobile station (B), the second mobile station (B) receiving the third message ( $M_3$ ) and verifying the validity of the information, if the information is valid the second mobile station (B) accepting the sharing of the shared encryption key K with the first mobile station (A).

7. A method according to claim 6, characterized in that the second mobile station (B) is authenticated by the first mobile station (A) selecting a prime number  $p$ , a generator  $a$  of a multiplicative group of integers modulo  $p$  when  $p \geq a \geq 2$  and a random secret  $x$  when  $p-2 \geq x \geq 1$ , constructing and sending to the second mobile station (B) the first message ( $M_1$ ) containing

$$a, p, a^x \bmod p,$$

- the second mobile station (B) receiving the first message ( $M_1$ ) and afterwards generating a secret  $y$  when  $p-2 \geq y \geq 1$  and computing a second shared key  $K_2 = (a^x)^y \bmod p$ , signing a concatenation of exponentials  $\{a^y, a^x\}$  and encrypting a result  $S_B\{a^y, a^x\}$  with the second shared key leading to  $E_K(S_B\{a^y, a^x\})$ , constructing and sending the second message ( $M_2$ ) to the first mobile station (A) containing

$$a^y \bmod p, cert_B, E_K(S_B\{a^y, a^x\}),$$

- certificate  $cert_B$  in the second message ( $M_2$ ) containing a signature verification key of the second mobile station (B), the exact contents of the certificate being of at least the following minimum

$$cert_B = (B, p_B, a, p, S_T\{B, p_B, a, p\}),$$

$p_B$  being a public signature verification key of the mobile station B and  $S_T$  a signature transformation of a trusted authority T whose public signature verification key is known in the first and second mobile stations (A, B),

- 20 the first mobile station (A) receiving the second message ( $M_2$ ) and afterwards computing a first shared encryption key  $(a^y)^x \bmod p = (a^x)^y \bmod p = K_1$ , checking the validity of the certificate  $cert_B$  the first mobile station (A), when the certificate  $cert_B$  is valid the encrypted part  $E_K(S_B\{a^y, a^x\})$  of the second message ( $M_2$ ) is decrypted to receive a signature  $S_B\{a^y, a^x\}$  and the signature  $S_B\{a^y, a^x\}$  is verified  
25 with a public signature verification key  $p_B$  of the second mobile station (B), if the signature  $S_B\{a^y, a^x\}$  is verified valid the first mobile station (A) accepts to share the shared encryption key  $K_1$  with the second mobile station (B),

the first mobile station (A) signing a concatenation of exponentials  $\{a^x, a^y\}$  and encrypting result  $S_A\{a^x, a^y\}$  with the first shared key  $K_1$  leading to  $E_K(S_A\{a^x, a^y\})$ ,

the first mobile station (A) constructing and sending the third message ( $M_3$ ) to the second mobile station (B) containing

$$cert_A, E_K(S_A\{a^x, a^y\}),$$

- 5  $cert_A$  including corresponding information with  $cert_B$  of the first mobile station (A), exact contents of the certificate  $cert_A$  being at least of the following minimum

$$cert_A = (B, p_A, a, p, S_T\{B, p_A, a, p\}),$$

$p_A$  being a public signature verification key of the first subscriber (A) and  $S_T$  a signature transformation of a trusted authority T whose public signature verification key is known by the first and second mobile stations (A, B),

- 10 the second mobile station (B) receiving the third message ( $M_3$ ) and verifying validity of the  $cert_A$ , decrypting  $E_A(S_A\{a^x, a^y\})$  and verifying validity of signature of  $S_A\{a^x, a^y\}$ , if all the signatures are valid the second mobile station (B) accepting sharing of the second shared encryption key  $K_2$  with the mobile station A.

- 15 8. A method according to any preceding claim 1 to 7, characterized in that the data is exchanged through user-to-user signalling.

9. A cellular communications system, where the first and second mobile stations (A, B) are wireless connected with via base stations, characterized in that it comprises

- 20 - a first mobile station (A, 9), that constructs and sends a first message ( $M_1$ ), receives and verifies the validity of a second message ( $M_2$ ) and when the information is verified valid accepts to share a shared encryption key K, constructs and sends a third message ( $M_3$ ),

- 25 - a second mobile station (B, 12), that receives the first message ( $M_1$ ) and constructs and sends the second message ( $M_2$ ), receives and verifies the validity of the third message ( $M_3$ ) and when the information is valid accepts to share the shared encryption key K with the first mobile station (A), and

- at least one mobile switching centre (10, 11).

10. A communications system according to claim 9, **characterized** in that it comprises two mobile switching centres (10, 11) connected together with ISDN.

11. A mobile station, **characterized** in that it comprises

- 5 - a processor (13) to perform operations needed to form and verify messages ( $M_1$ ,  $M_2$ ,  $M_3$ ), to implement authentication and key agreement procedures,
- a memory (14), where procedures and messages are stored with necessary parameters and variables,
- output means, on which commencement of extra secure communication is presented to a user of the mobile station,
- 10 - input means to enable validation of the extra secure communication,
- a transmitter/receiver (17) and an antenna (18) to transform information to radio waves from digital signals and vice versa.

12. A mobile station according to claim 11, **characterized** in that the output means comprises a display (15).

15 13. A mobile station according to claim 11, **characterized** in that the input means comprises a keyboard (16).

14. A mobile station according to claim 11, **characterized** in that it is designed to GSM standards.

20 15. A mobile station according to claim 11, **characterized** in that it is designed to UTMS standards.

**(57) Abstract**

The invention relates to a method to authenticate a mobile station B in a mobile network, so that the mobile station B is authenticated and an encryption key is agreed between mobile stations A and B using user data exchange during call setup. More specifically the mobile station B is authenticated by the mobile station A constructing and sending to the mobile station B a message  $M_1$ , the mobile station B receiving the message  $M_1$ , constructing and sending a message  $M_2$  to the mobile station A, the mobile station A receiving the message  $M_2$ , checking the validity of the information in the message  $M_2$ , if the information is verified valid the mobile station A accepting to share a shared encryption key K with mobile station B, the mobile station A constructing and sending the message  $M_3$  to the mobile station B, the mobile station B receiving the message  $M_3$  and verifying the validity of the information, if the information is valid the mobile station B accepting the sharing of the shared encryption key K with the mobile station A.



Figure 3



4  
13

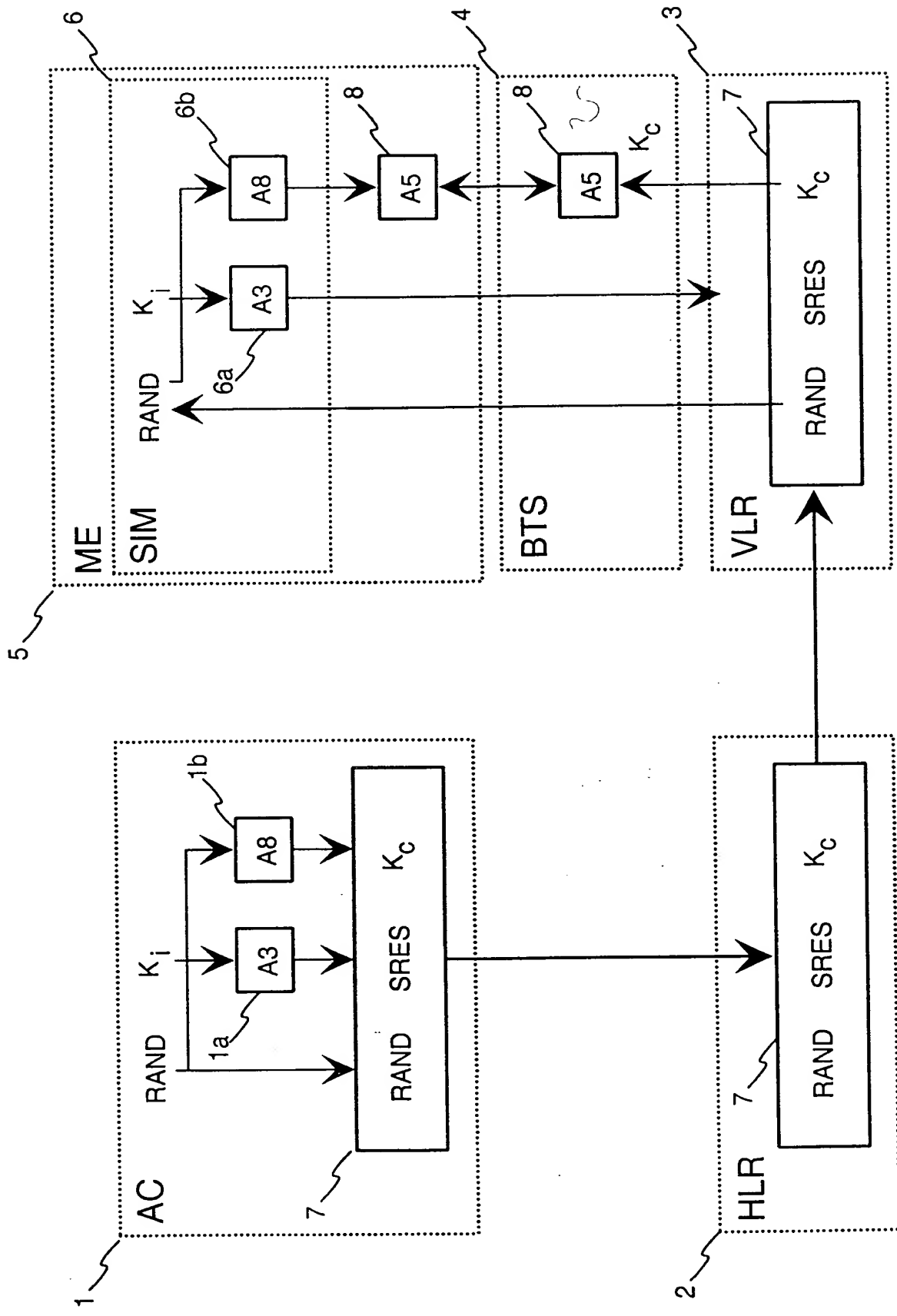


FIGURE 1

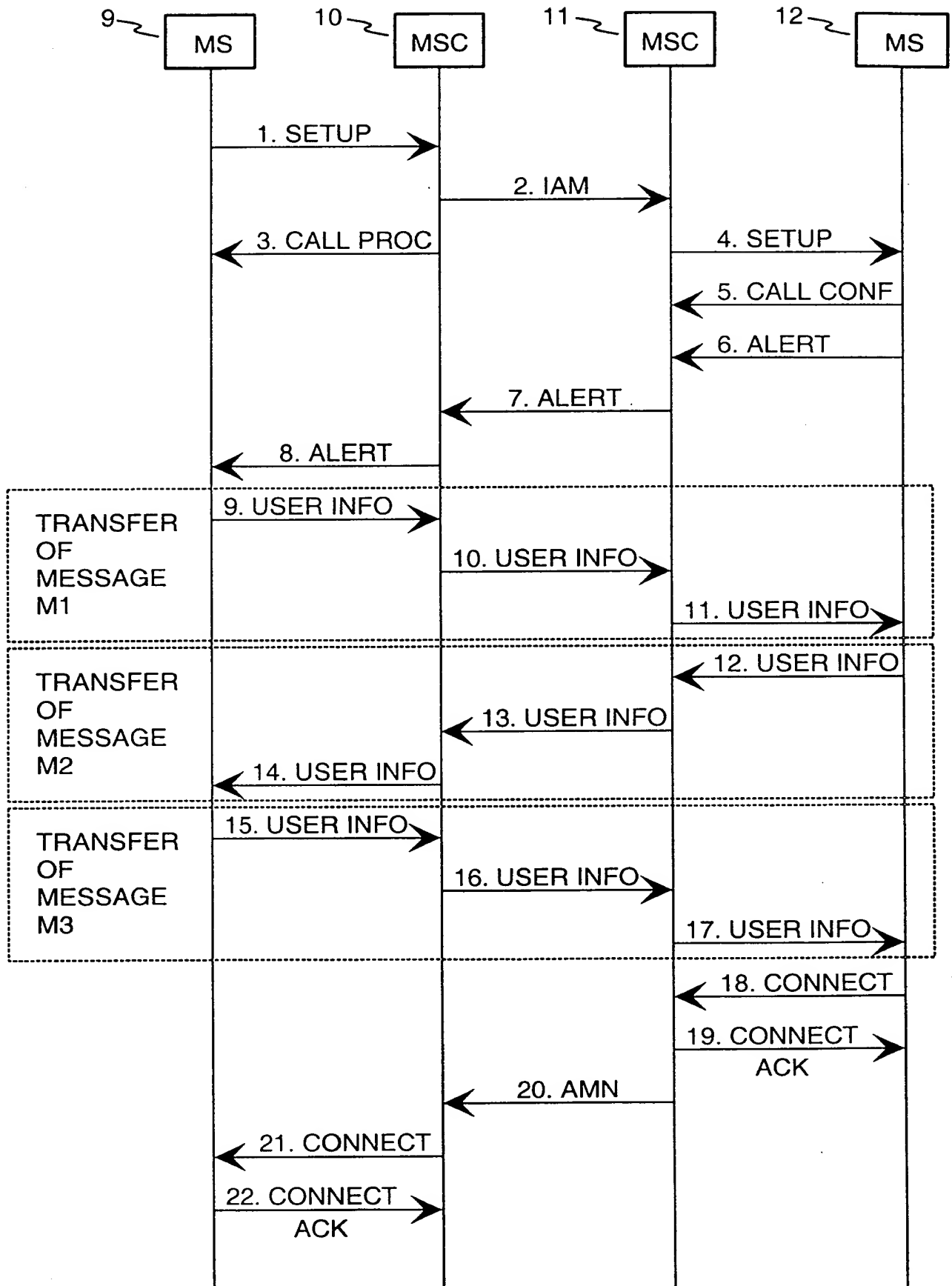


FIGURE 2

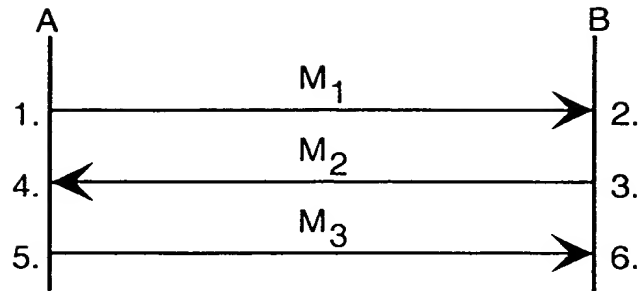


FIGURE 3

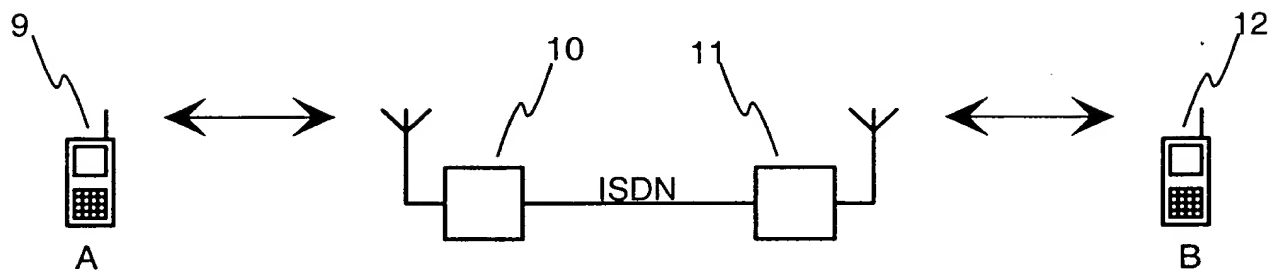


FIGURE 4

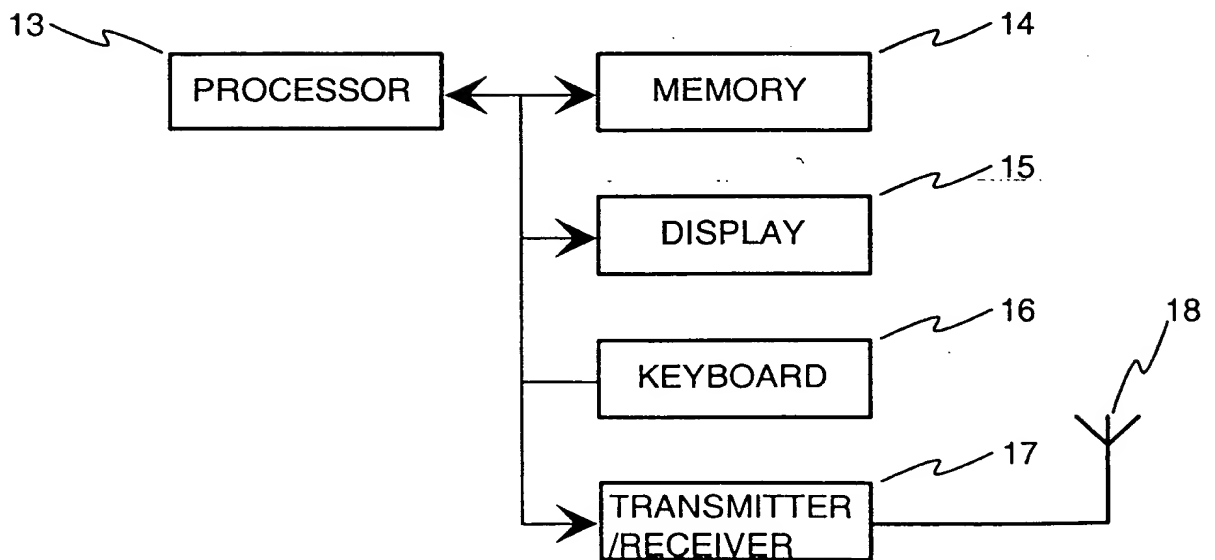


FIGURE 5